The preemptive stochastic resourceconstrained project scheduling problem: An efficient optimal solution procedure

Stefan Creemers<br>(December 5, 2016)

## Agenda

- Past work
- New approach
- What about the SRCPSP?
- Contribution


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- Past work

New approach
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## Past work: overview



> Creemers, Leus, Lambrecht (2010). Scheduling Markovian PERT networks to maximize the net present value, Operations
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1. Maximum-eNPV objective
2. No resources
3. Exponentially-distributed activity durations
4. Use of a SDP recursion to obtain the optimal policy

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1. Minimum-makespan objective
2. Renewable resources
3. General activity durations (PH approximation)
4. Use of an improved/modified SDP recursion

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## Past work

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## New approach



1. SDP recursion
2. Optimal solution
3. General activity durations
4. eNPV \& SRCPSP
5. UDCs to structure state space
6. Upper bound state space $=3^{n}$

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## Main bottleneck = memory!

## New approach



1. SDP recursion
2. Optimal solution
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5. UDCs to structure state space
6. Upper bound state space $=3^{n}$
 (D) Coses)
7. SDP recursion
8. Optimal solution
9. General activity durations
10. eNPV \& SRCPSP
11. No UDCs
12. Upper bound state space $=2^{n}$

Main bottleneck = memory!

## New approach: results

- Computational experiment to compare the old and the new approach with respect to:
- The number of instances solved
- The computation speed (CPU times)
- The average maximum number of states stored in memory
- We use a dataset with 30 projects for each:
- Number of activities ( $n$ between 10 \& 70)
- Order Strength (OS equal to 0.8, 0.6, and 0.4)


## New approach:

 number of instances solved| OLD |  |  |  |
| :---: | :---: | :---: | :---: |
| Number solved (out of 30) |  |  |  |
|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 30 | 30 | 30 |
| $n=20$ | 30 | 30 | 30 |
| $n=30$ | 30 | 30 | 30 |
| $n=40$ | 30 | 30 | 29 |
| $n=50$ | 30 | 30 | 16 |
| $n=60$ | 30 | 30 | 0 |
| $n=70$ | 30 | 29 | 0 |

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 number of instances solved| OLD |  |  |  |
| :---: | :---: | :---: | :---: |
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|  | OS = 0.8 | OS = 0.6 | OS $=0.4$ |
| $n=10$ | 30 | 30 | 30 |
| $n=20$ | 30 | 30 | 30 |
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| $n=40$ | 30 | 30 | 29 |
| $n=50$ | 30 | 30 | 16 |
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| NEW |  |  |  |
| :---: | :---: | :---: | :---: |
| Number solved (out of 30) |  |  |  |
|  | OS =0.8 | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 30 | 30 | 30 |
| $n=20$ | 30 | 30 | 30 |
| $n=30$ | 30 | 30 | 30 |
| $n=40$ | 30 | 30 | 30 |
| $n=50$ | 30 | 30 | 30 |
| $n=60$ | 30 | 30 | 30 |
| $n=70$ | 30 | 30 | 30 |

## New approach:

 average CPU time (sec)| OLD |  |  |  |
| :---: | :---: | :---: | :---: |
| Average CPU time (sec) |  |  |  |
|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 0.00 | 0.00 | 0.00 |
| $n=20$ | 0.00 | 0.00 | 0.00 |
| $n=30$ | 0.00 | 0.00 | 0.00 |
| $n=40$ | 0.00 | 0.00 | 41.1 |
| $n=50$ | 0.00 | 3.02 | 899 |
| $n=60$ | 0.00 | 39.4 | NA |
| $n=70$ | 0.00 | 365 | NA |

## New approach:

## average CPU time (sec)

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| :---: | :---: | :---: | :---: |
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|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 0.00 | 0.00 | 0.00 |
| $n=20$ | 0.00 | 0.00 | 0.00 |
| $n=30$ | 0.00 | 0.00 | 0.00 |
| $n=40$ | 0.00 | 0.00 | 41.1 |
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|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 0.00 | 0.00 | 0.00 |
| $n=20$ | 0.00 | 0.00 | 0.00 |
| $n=30$ | 0.00 | 0.00 | 0.00 |
| $n=40$ | 0.00 | 0.00 | 12.3 |
| $n=50$ | 0.00 | 0.00 | 270 |
| $n=60$ | 0.00 | 6.57 | 8960 |
| $n=70$ | 0.00 | 61.2 | 195691 |

On average, we improve computation times by a factor of 180 !

## New approach:

## average maximum number of states

| OLD |  |  |  |
| :---: | :---: | :---: | :---: |
| Average maximum \# states (x1000) |  |  |  |
|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $n=10$ | 0.00 | 0.00 | 0.00 |
| $n=20$ | 0.00 | 2.39 | 38.6 |
| $n=30$ | 0.00 | 24.8 | 934 |
| $n=40$ | 2.9 | 273 | 25413 |
| $n=50$ | 9.97 | 2155 | 315807 |
| $n=60$ | 37.9 | 21140 | NA |
| $n=70$ | 112 | 149925 | NA |

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## average maximum number of states

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| :---: | :---: | :---: | :---: |
| Average maximum \# states (x1000) |  |  |  |
|  | OS $=0.8$ | OS = 0.6 | OS $=0.4$ |
| $n=10$ | 0.00 | 0.00 | 0.00 |
| $n=20$ | 0.00 | 2.39 | 38.6 |
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| :---: | :---: | :---: | :---: |
| Average maximum \# states (x1000) |  |  |  |
|  | OS $=0.8$ | OS $=0.6$ | OS $=0.4$ |
| $\mathrm{n}=10$ | 0.00 | 0.00 | 0.00 |
| $\mathrm{n}=20$ | 0.00 | 0.00 | 0.00 |
| $\mathrm{n}=30$ | 0.00 | 0.00 | 2.87 |
| $\mathrm{n}=40$ | 0.00 | 1.28 | 30.4 |
| $\mathrm{n}=50$ | 0.00 | 4.87 | 210 |
| $\mathrm{n}=60$ | 0.00 | 20.2 | 1693 |
| $\mathrm{n}=70$ | 0.00 | 79.1 | 11006 |

On average, we reduce memory requirements by a factor of 364!

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- Past work
- New approach
- What about the SRCPSP?

Contribution

## SRCPSP results:

 computational performance|  | J30 |  |
| :--- | :---: | :---: |
|  | Old | New |
| Instances in set | 480 | 480 |
| Instances solved | 480 | 480 |
| Average CPU time (sec) | 0.48 | 0.02 |
| Average max \# states (x1000) | 176 | 1.99 |

## SRCPSP results:

 computational performance|  | J30 |  | J60 |  |
| :--- | :---: | :---: | :---: | :---: |
|  | Old | New | Old | New |
| Instances in set | 480 | 480 | 480 | 480 |
| Instances solved | 480 | 480 | 303 | 303 (480) |
| Average CPU time (sec) | 0.48 | 0.02 | 1591 | 81.6 |
| Average max \# states (x1000) | 176 | 1.99 | 374499 | 508 |

## SRCPSP results:

 computational performance|  | J30 |  | J60 |  |
| :--- | :---: | :---: | :---: | :---: |
|  | Old | New | Old | New |
| Instances in set | 480 | 480 | 480 | 480 |
| Instances solved | 480 | 480 | 303 | $303(480)$ |
| Average CPU time (sec) | 0.48 | 0.02 | 1591 | 81.6 |
| Average max \# states (x1000) | 176 | 1.99 | 374499 | 508 |

We are even able to solve 196 instances of the J90 dataset and 3 instances of the J120 dataset

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We improve the models of Creemers et al. (2010) and Creemers (2015) and obtain an increase in computational efficiency with factor 180 and a reduction of memory requirements with factor 364!

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We can use our model to find the optimal expected NPV for projects with up to 120 activities that have general activity durations!

## Contributions



We improve the models of Creemers et al. (2010) and Creemers (2015) and obtain an increase in computational efficiency with factor 180 and a reduction of memory requirements with factor 364 !


We can use our model to find the optimal expected NPV for projects with up to 120 activities that have general activity durations!


Our model can also be used to study the SRCPSP where the execution of activities is allowed to be interrupted (i.e., we can assess the value of splitting activities).
?

