Scheduling Markovian PERT networks with maximum-NPV objective

Stefan Creemers, Marc Lambrecht, Roel Leus

Department of Decision Sciences and Information Management

Katholieke Universiteit Leuven

April 28, 2008

◆□▶ ◆□▶ ◆注▶ ◆注▶ 注 のへで



Markov and Markov-regenerative PERT networks Kulkarni V.G. and Adlaka V.G. Operations Research (1986) Vol. 34(5) pp.769-781

(日) (同) (三) (三)

3

 Problem Description
 Markov PERT Networks

 Model
 NPV-Objective in Stochastic Project Networks

 Results
 Contribution

- Markov and Markov-regenerative PERT networks Kulkarni V.G. and Adlaka V.G. Operations Research (1986) Vol. 34(5) pp.769-781
 - PERT networks with independent exponentially distributed activity durations
 - Project execution is a Continuous Time Markov Chain with a single absorbing state (i.e. project completion)
 - Early-start policy is optimal

(日)

NPV is a nonregular measure of performance, starting activities as soon as possible is not necessarily optimal

Extensive body of literature exists on the deterministic case:

- A.H. Russell (1970)
- R.C. Grinold (1972)
- S. Elmaghraby and W. Herroelen (1990)
- R.H. Möhring, A.S. Schulz, F. Stork and M. Uetz (2001)
- C. Schwindt and J. Zimmermann (2001)

< ロ > < 同 > < 回 > < 回 > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □





Activity Delay in Stochastic Project Networks Buss A.H. and Rosenblatt M.J. Operations Research (1997) Vol. 45(1) pp. 126-139

Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

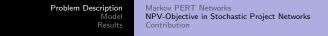
(日) (同) (三) (三)

3



- Activity Delay in Stochastic Project Networks Buss A.H. and Rosenblatt M.J. Operations Research (1997) Vol. 45(1) pp. 126-139
 - Algorithms to determine delays at the onset of the project (i.e. static decisions)
 - Early-start policy after delay
 - Performance limited to 25-activity networks

イロト イポト イヨト イヨト



 Scheduling projects with stochastic activity duration to maximize EPV
 Tilson V., Sobel M.J. and Szmerekovsky J.G.
 Submitted Working Paper (2006)

イロン 不同 とくほう イロン

3



 Scheduling projects with stochastic activity duration to maximize EPV
 Tilson V., Sobel M.J. and Szmerekovsky J.G.
 Submitted Working Paper (2006)

- Optimization over the set of policies that start activities at the end of other activities (dynamic)
- Process is a Continuous Time Markov Decision Chain
- Performance limited to 25-activity networks

-



Our contribution:

Significant improvement of performance compared to existing models:

- CPU-time reduction up to factor 15
- Memory requirement reduction up to factor 360 (largest statespace analyzed: 867,589,281 states)

< ロ > < 同 > < 回 > < 回 > < 回 > <

Setting:

- Stochastic activity durations (exponentially distributed)
- Expected NPV-objective: incurred cash flow c_i at the start of activity i
- Optimization over all policies that start activities at the end of other activities
- No resources

Model outline:

- Definition of the statespace
- Dynamic program to obtain optimal NPV

- 4 同 6 4 日 6 4 日 6

Preliminary Concepts

Status of activity *i* at time *t*:

- not started $\Omega_i(t) = 0$
- in progress $\Omega_i(t) = 1$
- finished $\Omega_i(t) = 2$

 $\Omega(t) = (\Omega_0(t), \Omega_1(t), \dots, \Omega_n(t))$ defines the state of the system

(人間) ト く ヨ ト く ヨ ト

Preliminary Concepts

Status of activity *i* at time *t*:

- not started $\Omega_i(t) = 0$
- in progress $\Omega_i(t) = 1$
- finished $\Omega_i(t) = 2$

 $\Omega(t) = (\Omega_0(t), \Omega_1(t), \dots, \Omega_n(t))$ defines the state of the system

Size of statespace Q has upper bound $|Q| = 3^n$

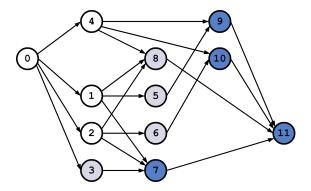
Most of these states do not satisfy precedence constraints \Rightarrow Strict and clear definition of the statespace is essential

< ロ > < 同 > < 回 > < 回 > < □ > <

Problem Description Model Results

Definition of the Statespace Dynamic Program

Example of a Feasible State



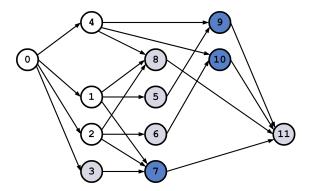
Feasible state $\Omega = (2, 2, 2, 1, 2, 1, 1, 0, 1, 0, 0, 0)$

イロト イボト イヨト イヨト

Problem Description Model Results

Definition of the Statespace Dynamic Program

Example of an Infeasible State



Infeasible state $\Omega = (2, 2, 2, 1, 2, 1, 1, 0, 1, 0, 0, 1)$

イロト イボト イヨト イヨト

The UDC-concept

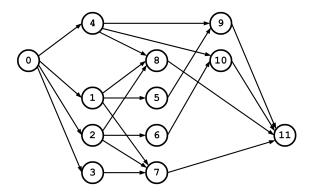
A UDC is an inclusion-maximal set of activities that can be executed in parallel at a given moment in time

Traditionally used in AoA representation, we apply the concept in AoN representation

(4月) (4日) (4日)

The UDC-concept

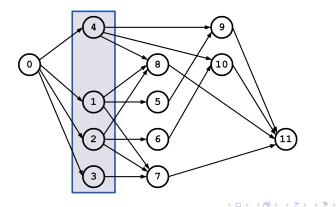
A UDC is a set of all activities that can be executed in parallel at a given moment in time



→ □ ▶ → 臣 ▶ → 臣 ▶

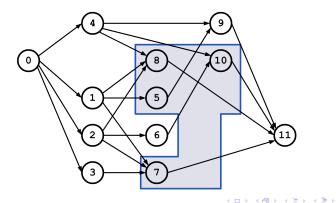
The UDC-concept

A UDC is a set of all activities that can be executed in parallel at a given moment in time



The UDC-concept

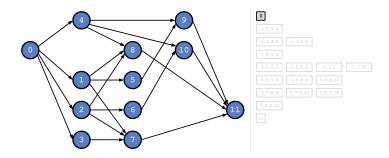
A UDC is a set of all activities that can be executed in parallel at a given moment in time



Problem Description Model Results

Definition of the Statespace Dynamic Program

The UDC-network



State $\Omega = (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$

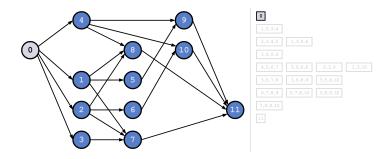
Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

- 4 同 6 4 日 6 4 日 6

Problem Description Model Results

Definition of the Statespace Dynamic Program

The UDC-network

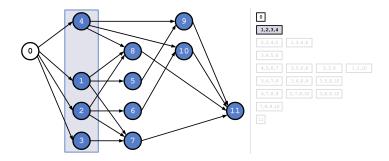


State $\Omega = (1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$

Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

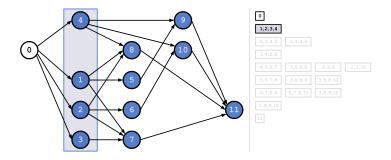
The UDC-network



State $\Omega = (2, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

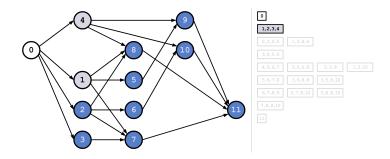
The UDC-network



State $\Omega = (2, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$ Lemma 1. Each feasible state is assigned to a single UDC

(日) (同) (三) (三)

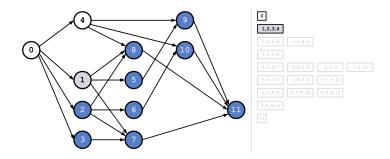
The UDC-network



State $\Omega = (2, 1, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0)$

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

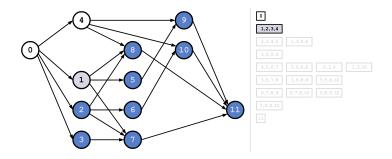
The UDC-network



State $\Omega = (2, 1, 0, 0, 2, 0, 0, 0, 0, 0, 0, 0)$

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

The UDC-network

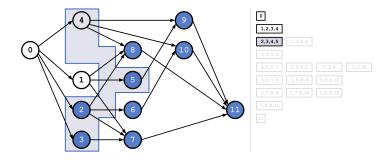


State $\Omega = (2, 1, 0, 0, 2, 0, 0, 0, 0, 0, 0, 0)$

Lemma 2. If at least one new activity becomes eligible then the system moves to a different UDC

(日) (同) (三) (三)

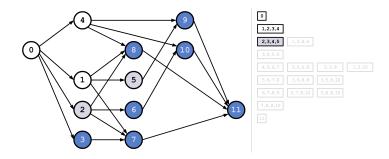
The UDC-network



State $\Omega = (2, 2, 0, 0, 2, 0, 0, 0, 0, 0, 0, 0)$

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

The UDC-network



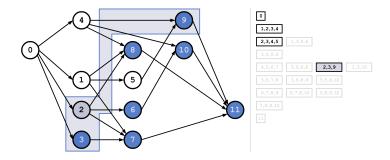
State $\Omega = (2, 2, 1, 0, 2, 1, 0, 0, 0, 0, 0, 0)$

- 4 回 2 - 4 □ 2 - 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □ 0 − 4 □

Problem Description Model Results

Definition of the Statespace Dynamic Program

The UDC-network



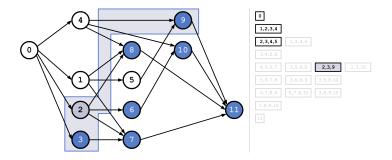
State $\Omega = (2, 2, 1, 0, 2, 2, 0, 0, 0, 0, 0, 0)$

Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

Problem Description Model Results

Definition of the Statespace Dynamic Program

The UDC-network

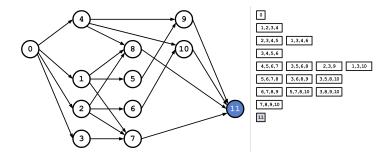


State $\Omega = (2, 2, 1, 0, 2, 2, 0, 0, 0, 0, 0, 0)$

Lemma 3. Inter-UDC-transitions can only lead from lower- to higher-ranked UDCs

(日) (同) (三) (三)

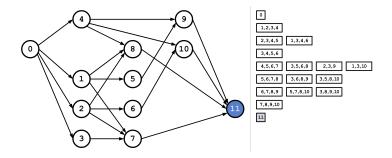
The UDC-network



State $\Omega = (2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 0)$

- 4 同 6 4 日 6 4 日 6

The UDC-network



State $\Omega = (2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 0)$

Observation 1. Note that the assignment of states to UDCs establishes a partition of Q

Algorithm Global algorithmic structure Generate the UDC-network Let NPV at state $\Omega = (2, 2, \dots, 2, 0)$ equal c_n For all UDCs in decreasing rank Allocate storage for all states in the UDC For all states Determine optimal decision and compute NPV (SDP-recursion) **End For** For all UDCs that are linked to the current UDC Reduce the number of incoming links If there are no more incoming links Free storage occupied by the UDC End If **End For** End For

- 4 同 🕨 - 4 目 🕨 - 4 目

Lemma 4. For an arbitrary UDC, in any state, the backward recursion only needs value-function lookups within higher ranked UDCs or within the same UDC for states which have already been evaluated.

(日)

Problem Description Model Results	Comparison with existing models Computational Results Memory Efficiency Time for Questions
--	---

	Tilson et al.	Creemers et al.
Configuration	Intel Pentium IV	AMD Athlon 64
Clock Speed	2.8GHz	1.8GHz
PCMark* (CPU)	3,646	2,602
RAM	512MB	2,048MB
CPU time	210 sec	14 sec
Max statespace	600,000	268, 435, 456
		(867, 589, 281)

CPU time reduction: factor 15 (uncorrected) Memory reduction: factor 360 (corrected)

*CPU benchmarking tool: http://www.futuremark.com/

▲ □ ▶ ▲ □ ▶ ▲ □ ▶

Problem Description Model Results	Comparison with existing models Computational Results Memory Efficiency Time for Questions
--	---

N	Ns			Average statespace size			
	<i>OS</i> = 0.8	OS = 0.6	OS = 0.4	OS = 0.8	OS = 0.6	OS = 0.4	
10	30	30	30	71	206	695	
20	30	30	30	484	4,006	55,016	
30	30	30	30	1,995	49, 388	1,560,364	
40	30	30	29	7,860	534,014	47,072,515	
50	30	30	4	26,667	4, 346, 215	526,020,237	
60	30	30	0	92,003	216,027,815		
70	30	22	0	286,831	216,027,815		
80	30	5	0	829, 741	758, 644, 207		
90	30	0	0	2,596,419			
100	30	0	0	6,868,100			
110	30	0	0	24, 235, 588			
120							

N	Avg	g CPU Time N	IPV	Ma	ax CPU Time NP\	/	
	<i>OS</i> = 0.8	OS = 0.6	OS = 0.4	OS = 0.8	OS = 0.6	OS = 0.4	
10	0.00	0.00	0.00	0.00	0.00	0.03	
20	0.00	0.03	0.90	0.00	0.08	3.77	
30	0.01	0.64	52.98	0.02	1.78	326.16	
40	0.06	13.29	4, 273	0.11	51.72	19,916	
50	0.27	171.56	99, 216	0.66	849.42	132, 984	
60	1.28	0.00		5.30	0.00		
70	5.37	33, 203		15.52	114, 424		
80	19.13	124,831		73.09	145, 922		
90	86.86			538			
100	301			1,626			
110	1,774			19, 571			
120							
							_
						◆ 差 ▶ < 差 ▶	

Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

Problem Description Model Results	Comparison with existing models Computational Results Memory Efficiency Time for Questions
--	---

N	Av	verage statespace	size	Maximum statespace size			
	OS = 0.8	OS = 0.6	OS = 0.4	OS = 0.8	OS = 0.6	OS = 0.4	
10	71	206	695	105	333	2,361	
20	484	4,006	55,016	953	7,673	153, 441	
30	1,995	49, 388	1,560,364	3, 233	84,837	5,966,721	
40	7,860	534,014	47,072,515	11,945	1,543,113	146, 560, 473	
50	26,667	4, 346, 215	526,020,237	53, 481	13, 893, 741	737,047,953	
60	92,003	42, 278, 506		236,889	165, 102, 585		
70	286,831	216,027,815		605,649	426, 644, 253		
80	829,741	758,644,207		2,278,353	867, 589, 281		
90	2, 596, 419			9, 322, 153			
100	6,868,100			22,963,321			
110	24, 235, 588			117, 261, 489			
120							

Ν	Maxi	mum statespace	use	Aver	age statespace us	e
	OS = 0.8	OS = 0.6	OS = 0.4	<i>OS</i> = 0.8	OS = 0.6	OS = 0.4
10	0.41	0.55	0.63	0.25	0.37	0.44
20	0.38	0.49	0.62	0.22	0.27	0.38
30	0.30	0.44	0.55	0.15	0.24	0.30
40	0.31	0.46	0.52	0.15	0.28	0.29
50	0.33	0.46	0.28	0.16	0.24	0.17
60	0.37	0.49		0.16	0.33	
70	0.34	0.40		0.16	0.19	
80	0.30	0.25		0.13	0.11	
90	0.35			0.16		
100	0.39			0.17		
110	0.42			0.19		
120						

Stefan Creemers, Marc Lambrecht, Roel Leus Scheduling Markovian PERT networks with maximum-NPV objective

Problem Description Model Results	Comparison with existing models Computational Results Memory Efficiency Time for Questions
--	---



◆□▶ ◆圖▶ ◆臣▶ ◆臣▶ 三臣…