



KATHOLIEKE UNIVERSITEIT  
**LEUVEN**



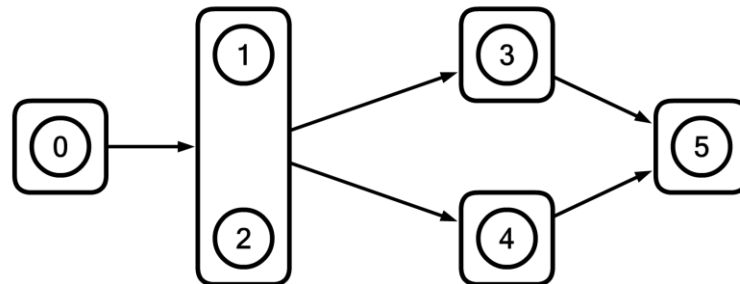
# Project Scheduling with Alternative Technologies and Stochastic Activity Durations

PMS Tours, April 2010

Stefan Creemers  
Roel Leus  
Bert De Reyck

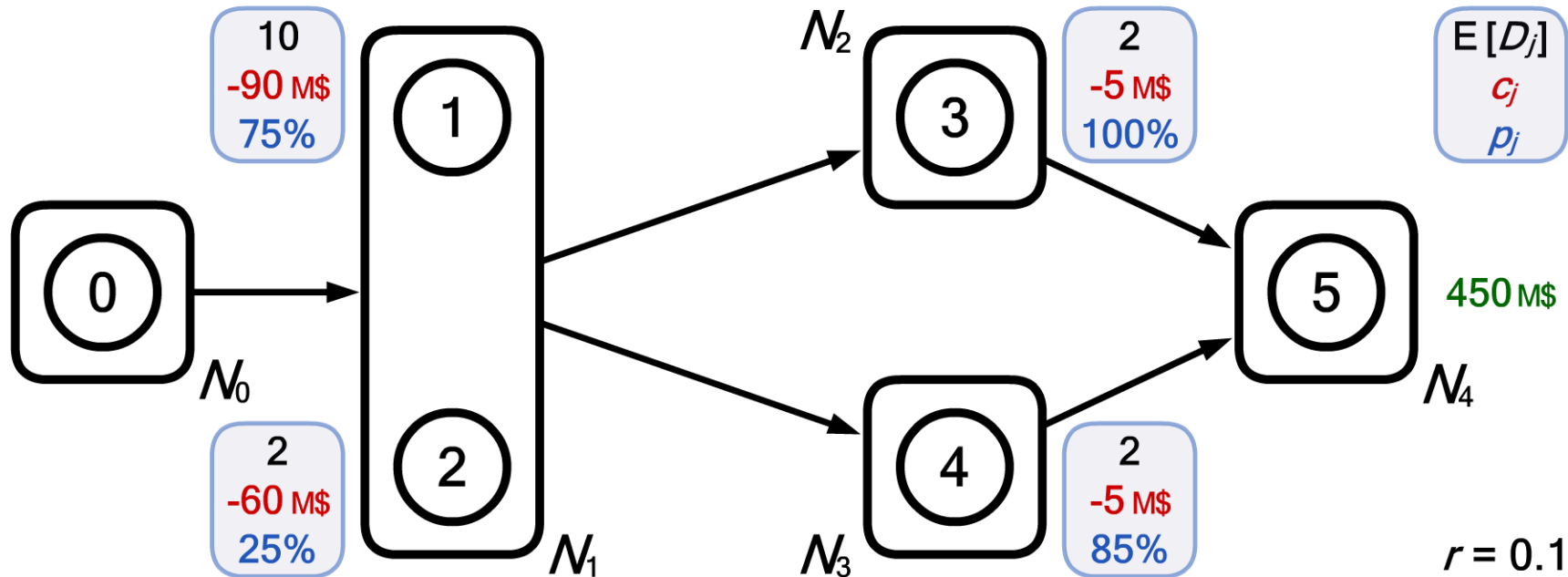
# Introduction

- Goal = maximize NPV of projects in which:
  - Activities can fail
  - Activities that pursue the same result may be grouped in “modules”
  - Each module needs to be successful for the project to succeed
  - A module is successful if at least one of its activities succeeds



- This is common in R&D (especially in NPD) but also in other sectors: pharmaceuticals, software development, ...

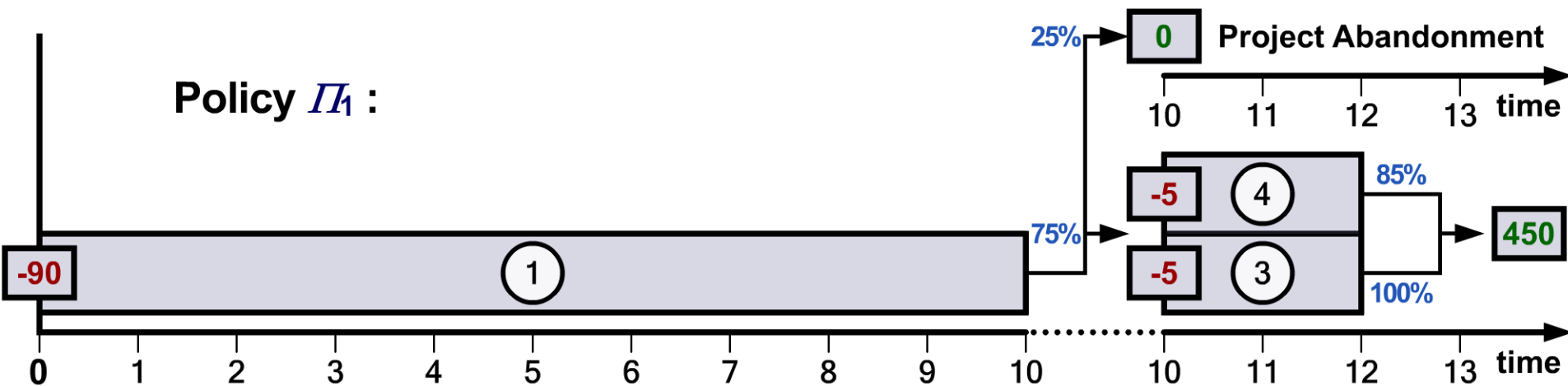
# Definitions



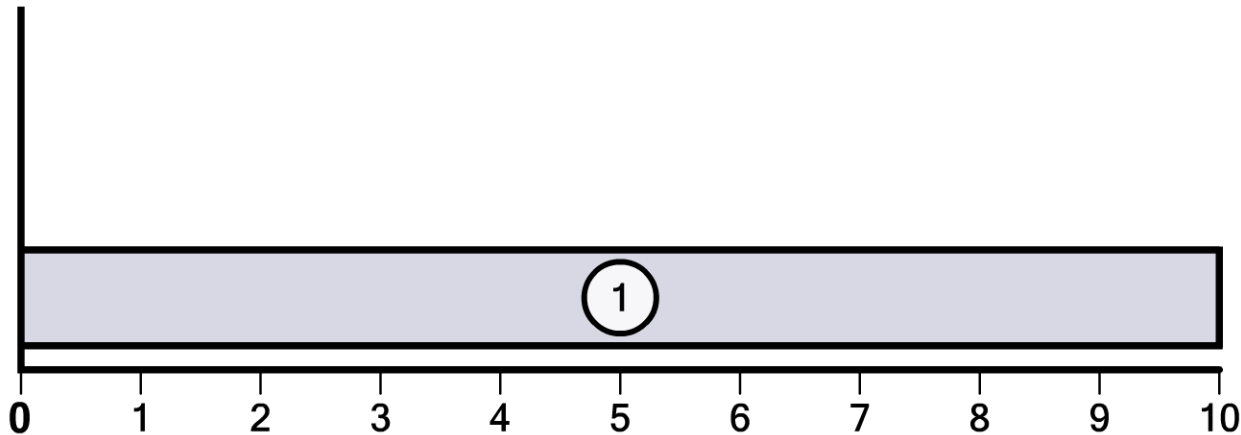
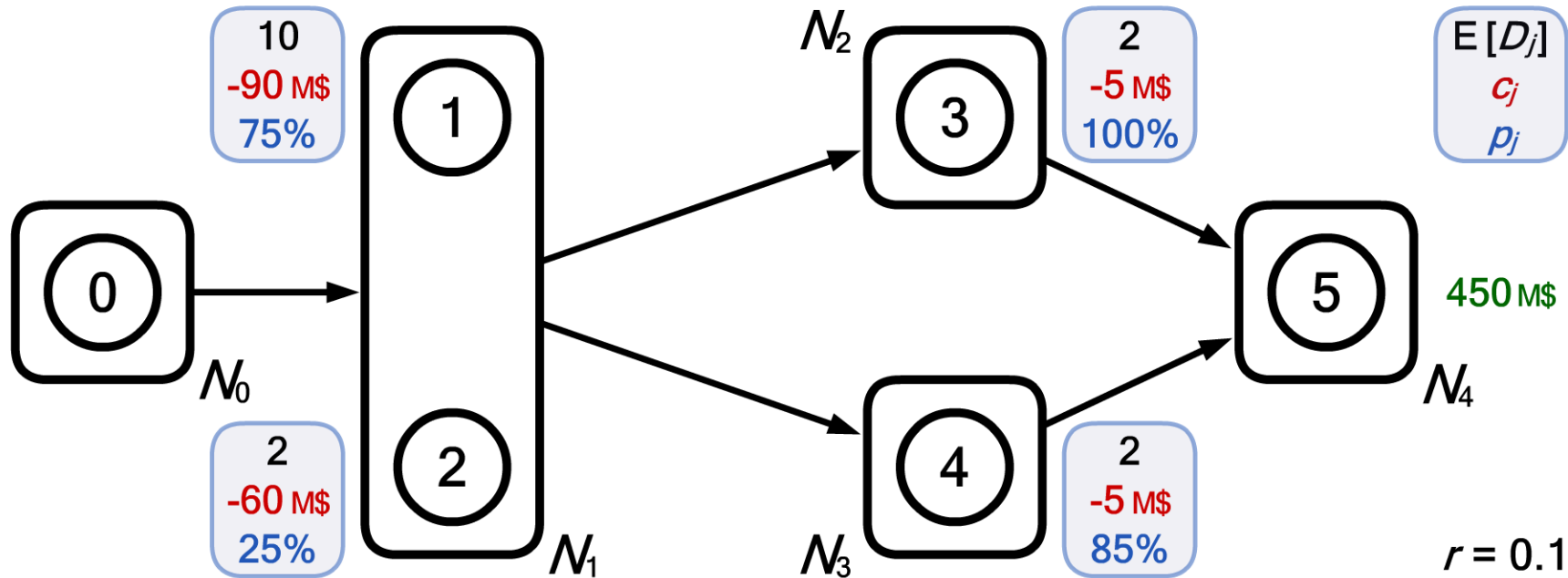
- Project network with  $n$  activities (activity = on the node)
- Stochastic activity durations: expected duration  $E[D_j]$  of activity  $j$
- Expected-NPV objective: cash flow  $c_j$  is incurred at the start of activity  $j$
- End-of-project payoff  $C$  obtained upon overall project success
- Failures: each activity  $j$  has a probability of technical success  $p_j$
- Time value of money  $\Rightarrow$  discount rate  $r$
- $m$  modules  $N_i$

# Deterministic durations

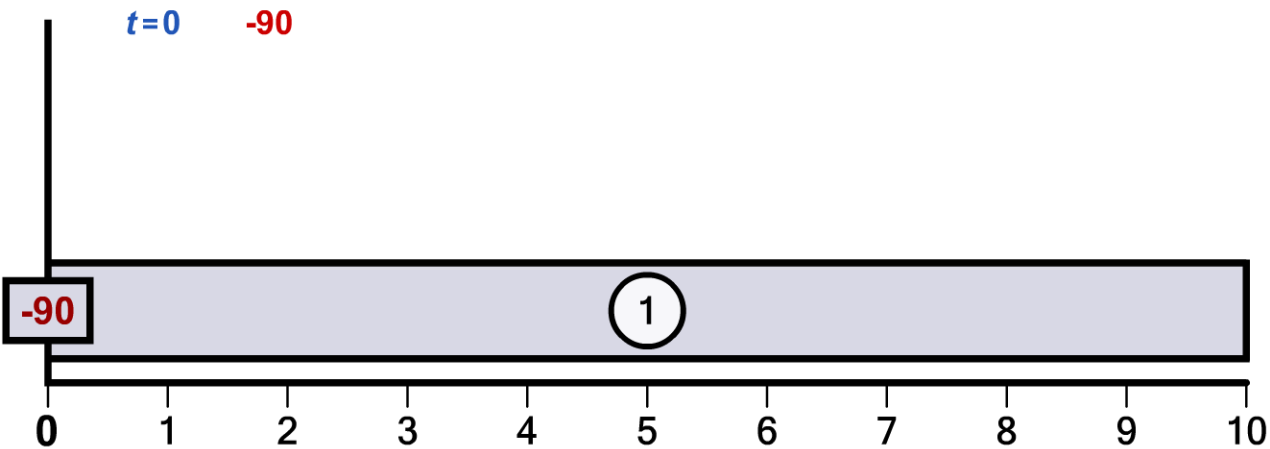
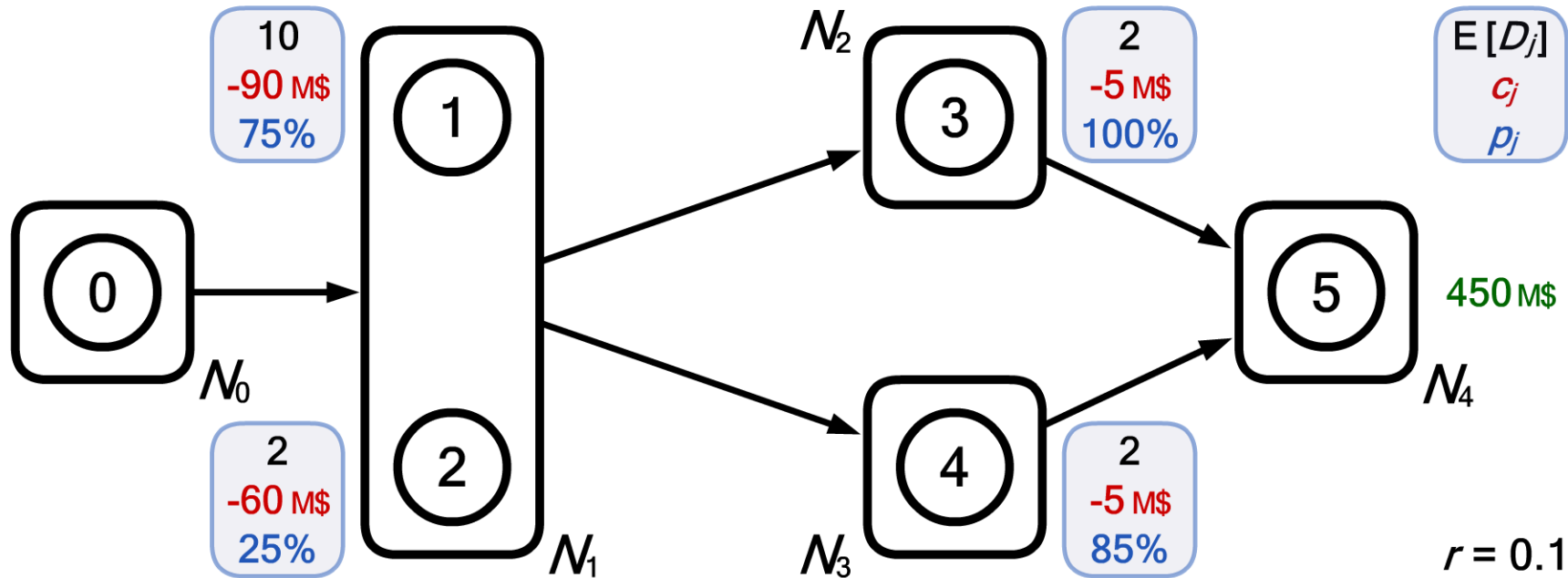
A solution is not a schedule but rather a scheduling policy (even with deterministic durations)



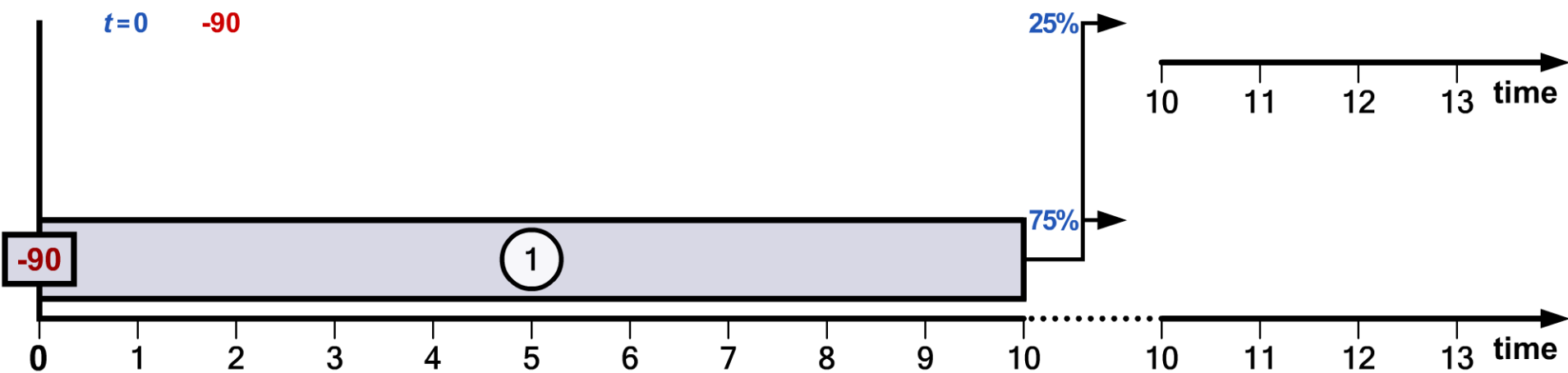
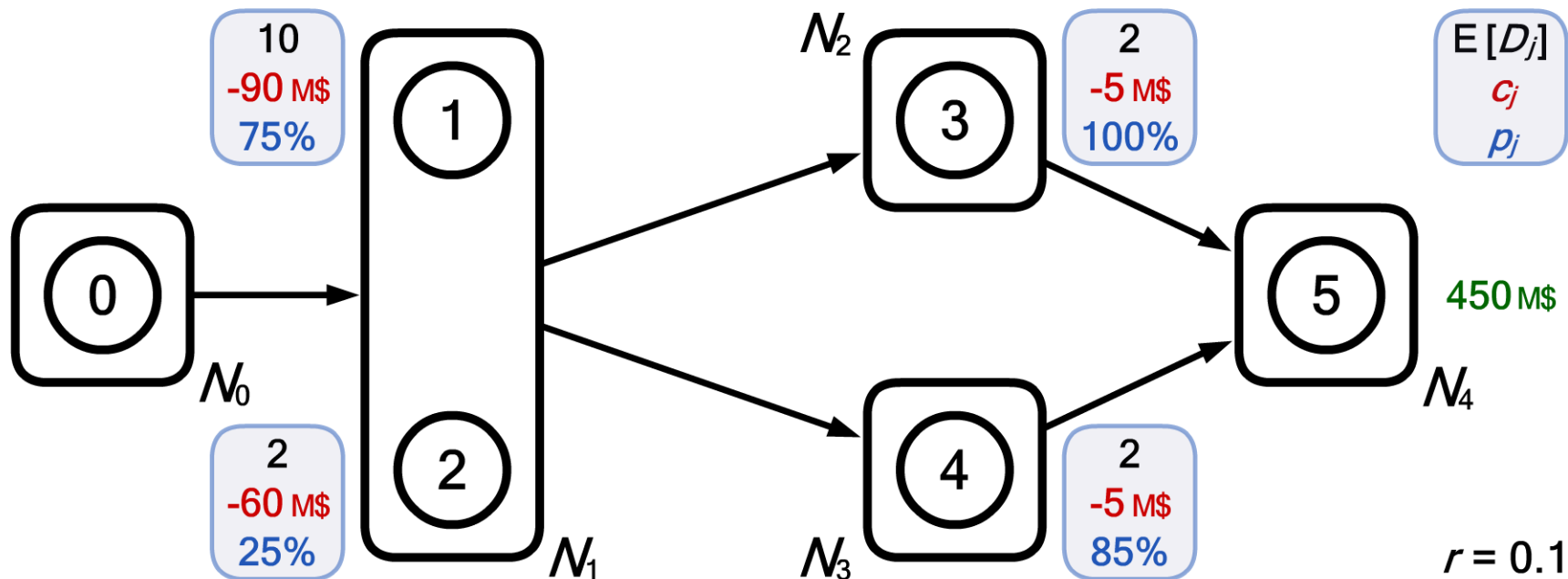
# Deterministic durations



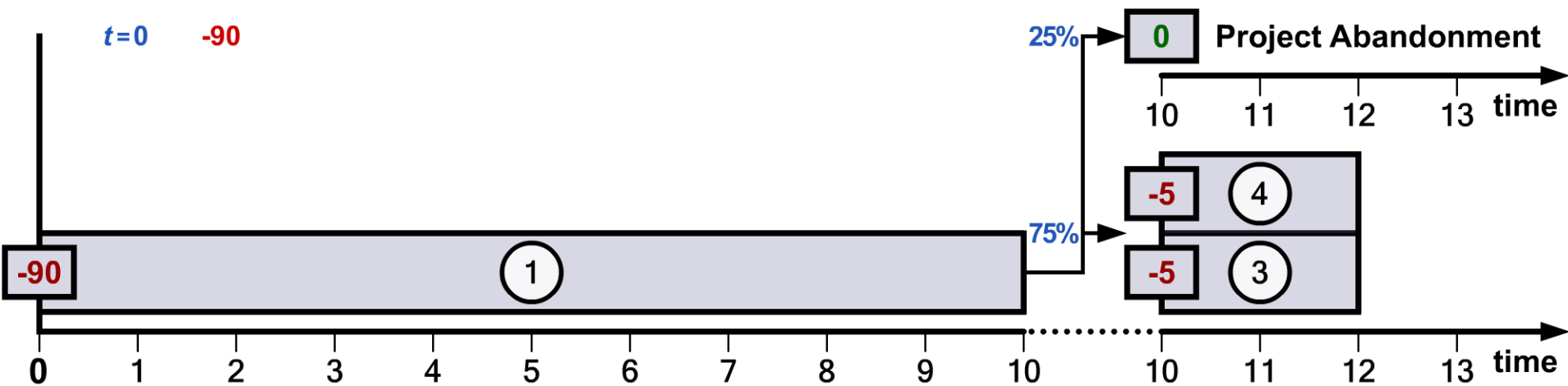
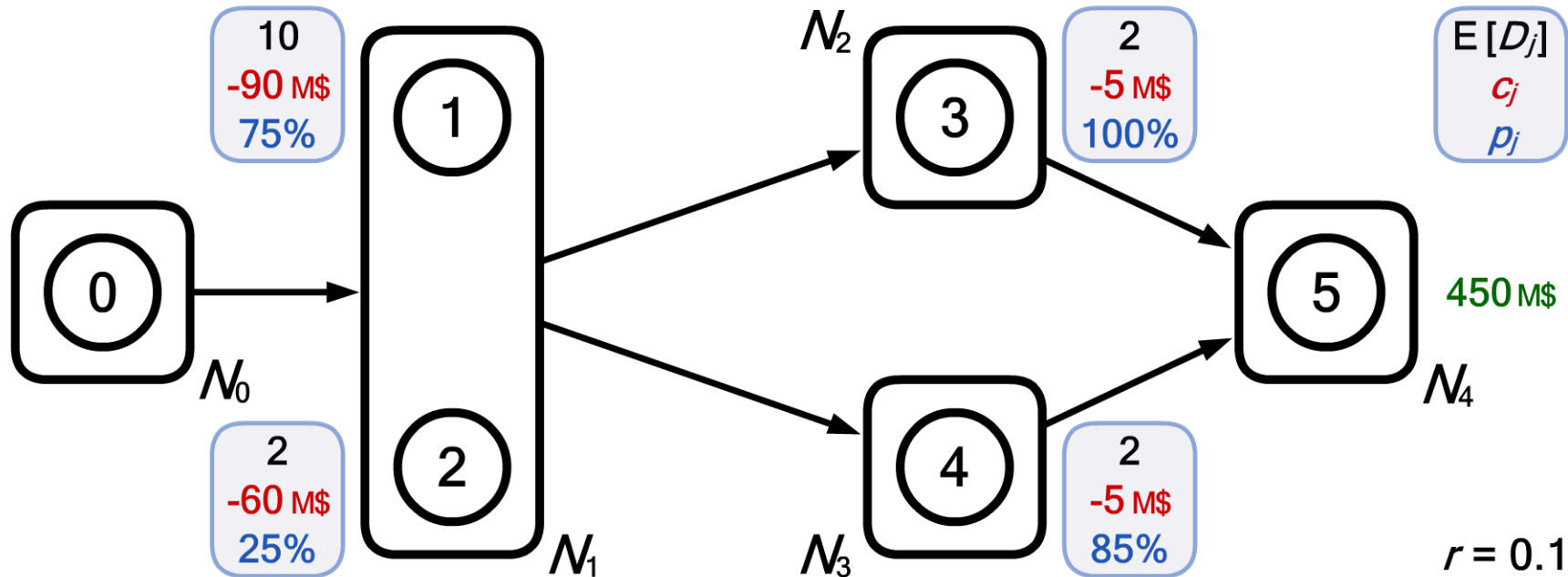
# Deterministic durations



# Deterministic durations

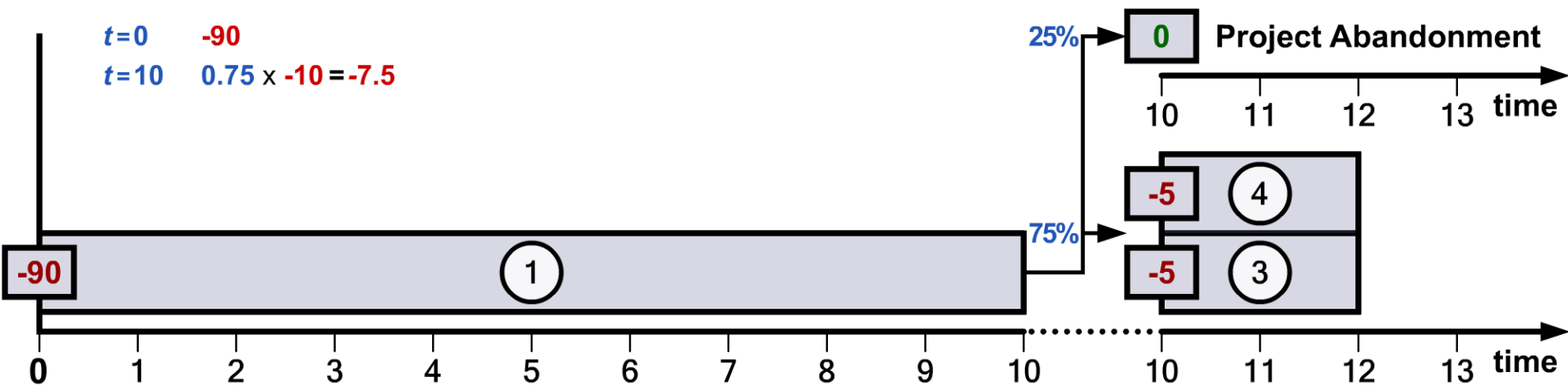
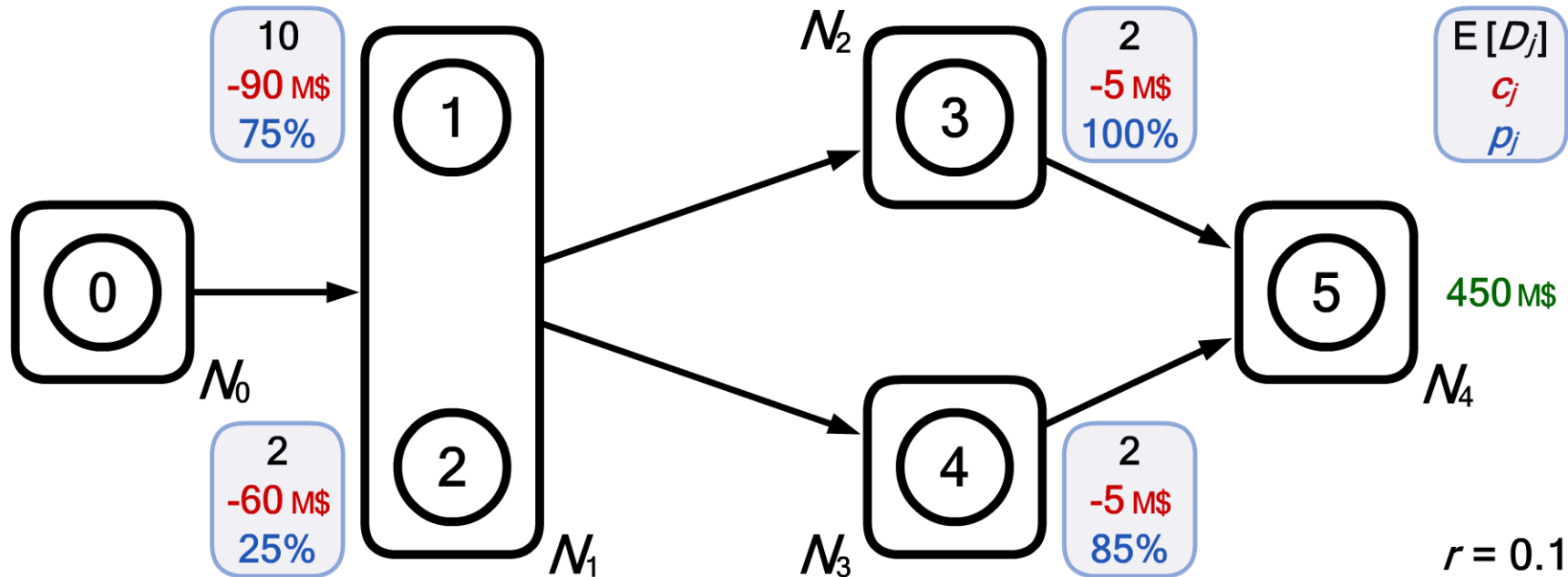


# Deterministic durations

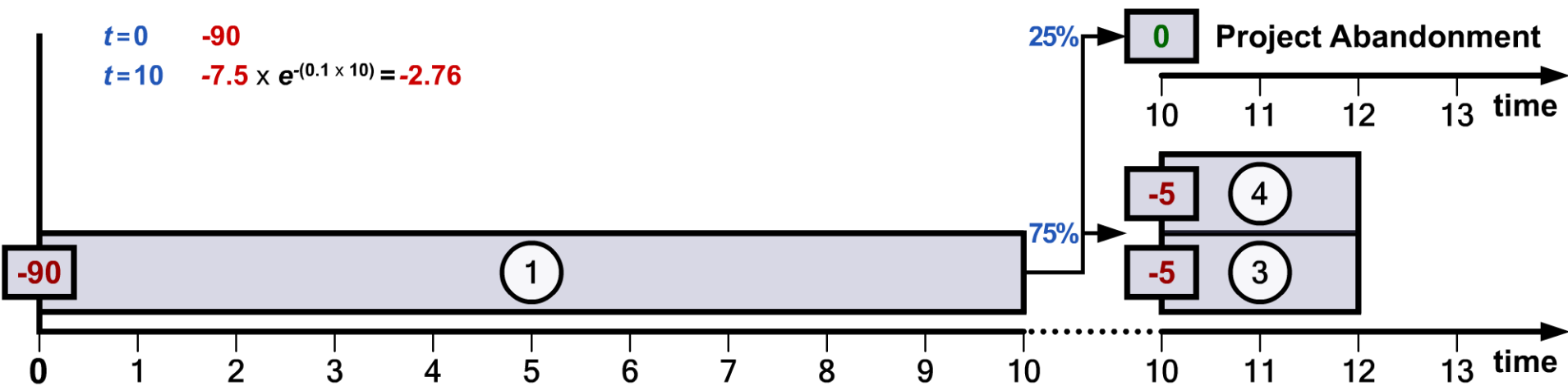
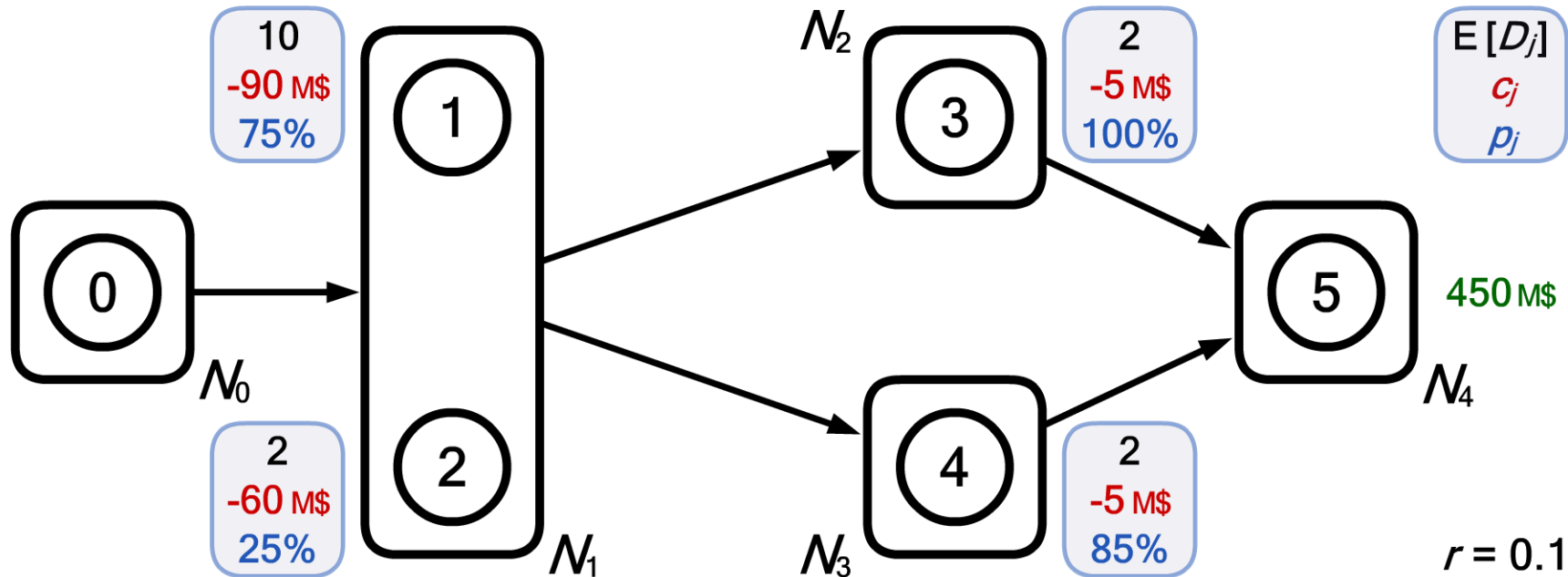




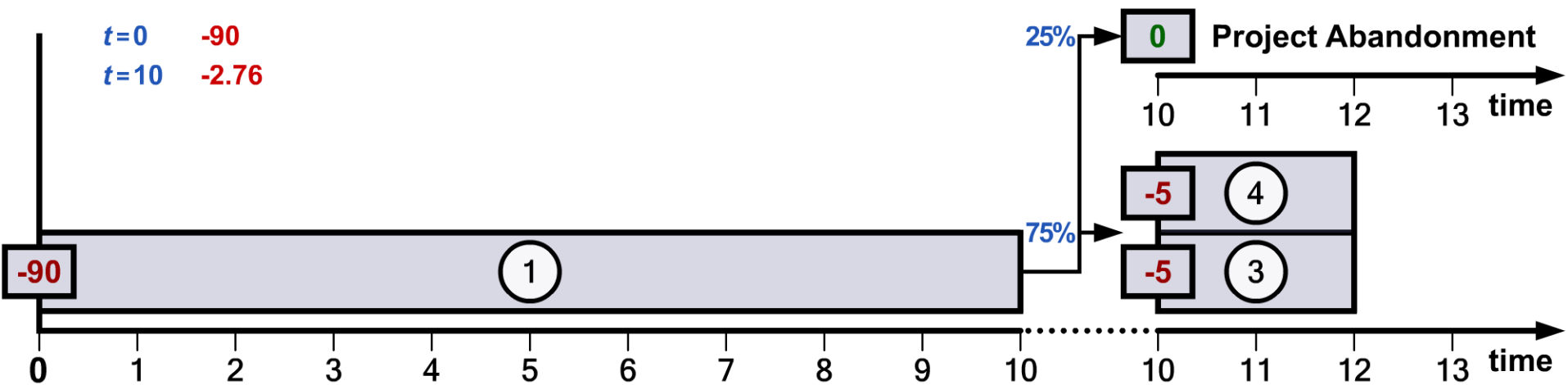
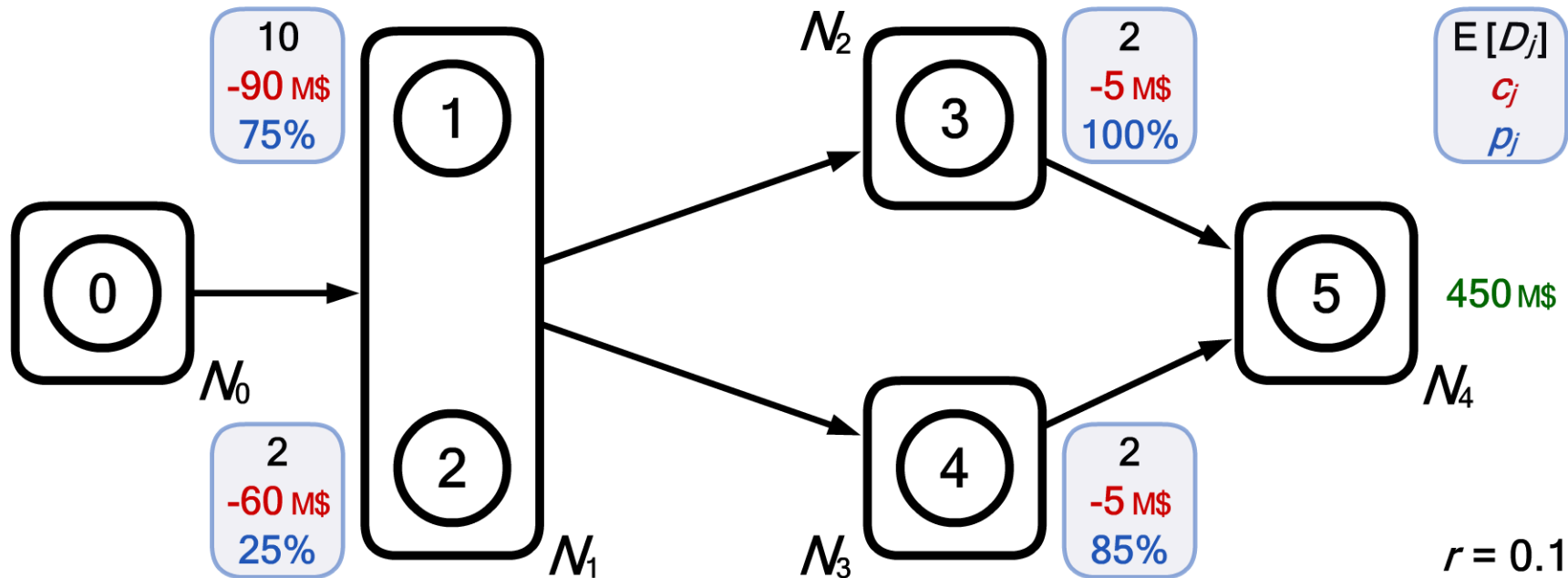
# Deterministic durations



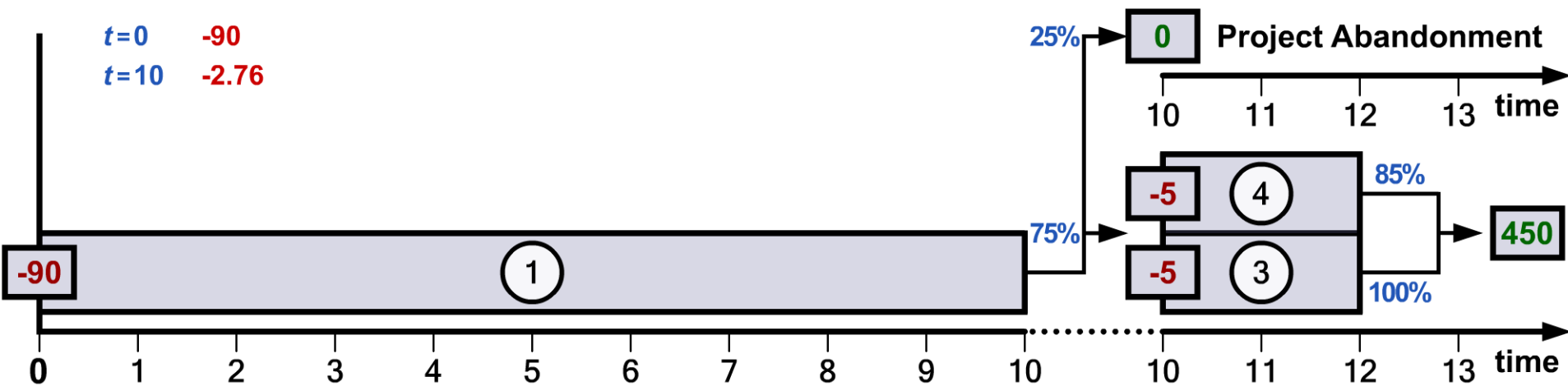
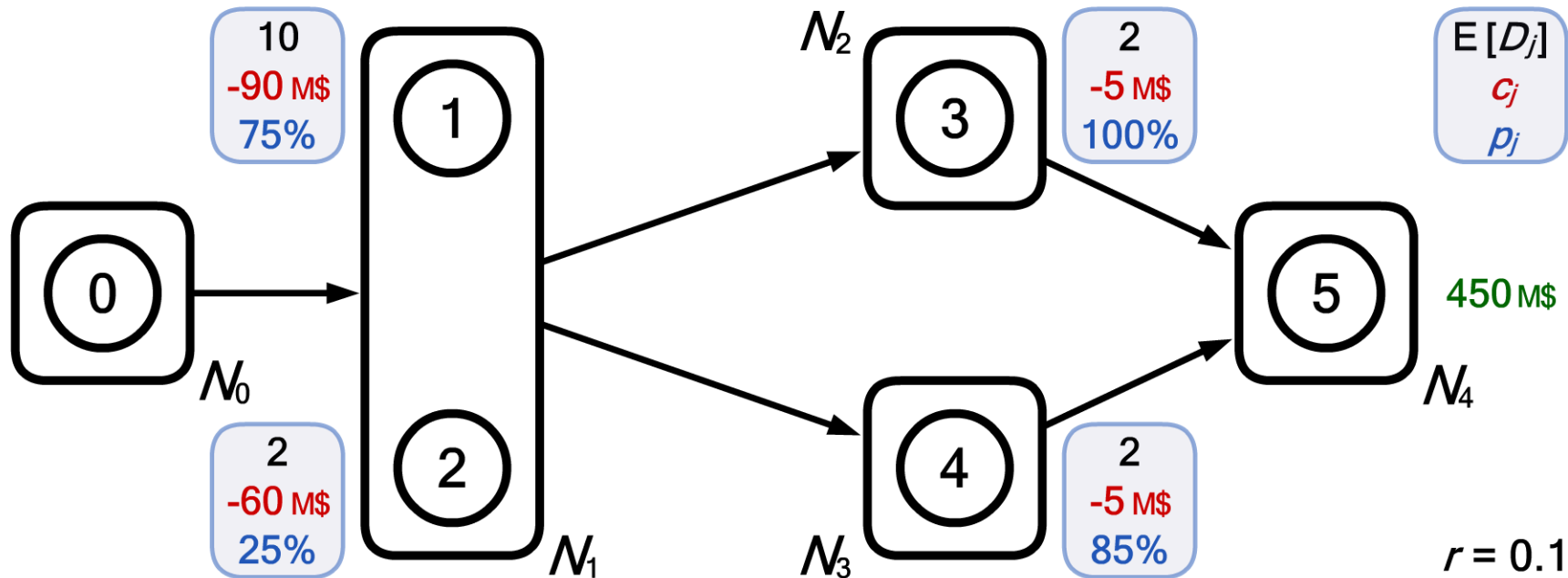
# Deterministic durations



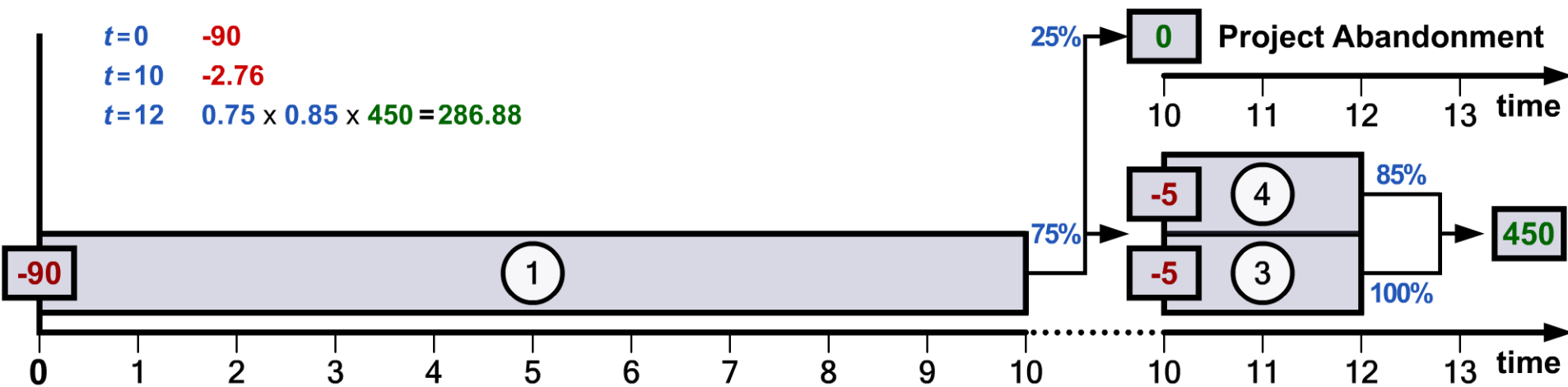
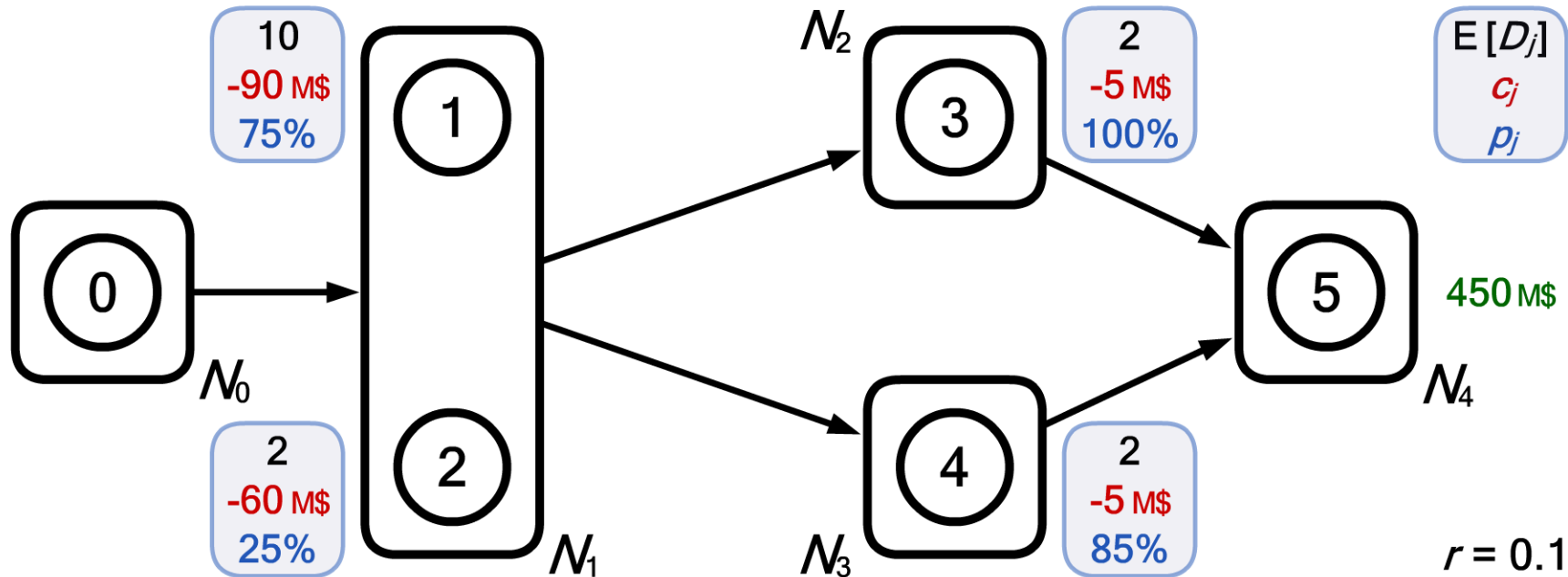
# Deterministic durations



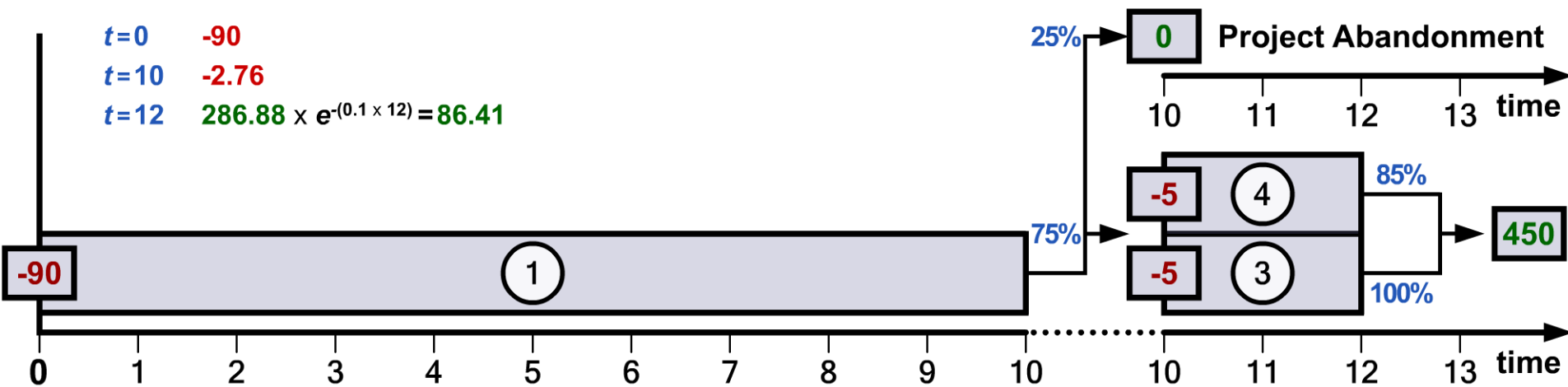
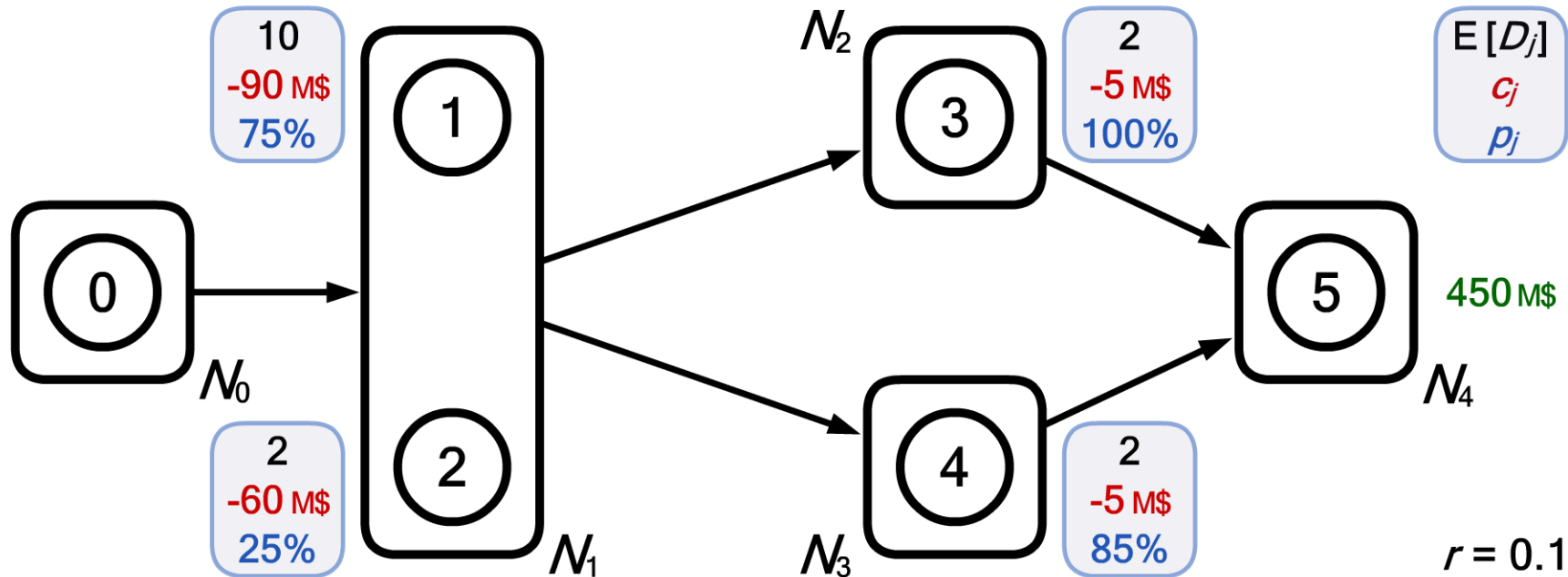
# Deterministic durations



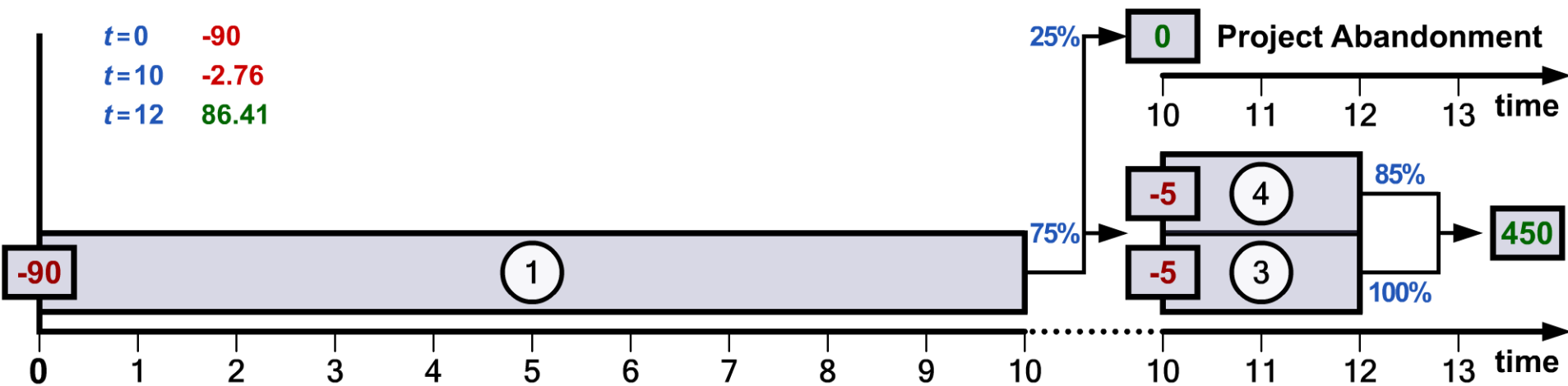
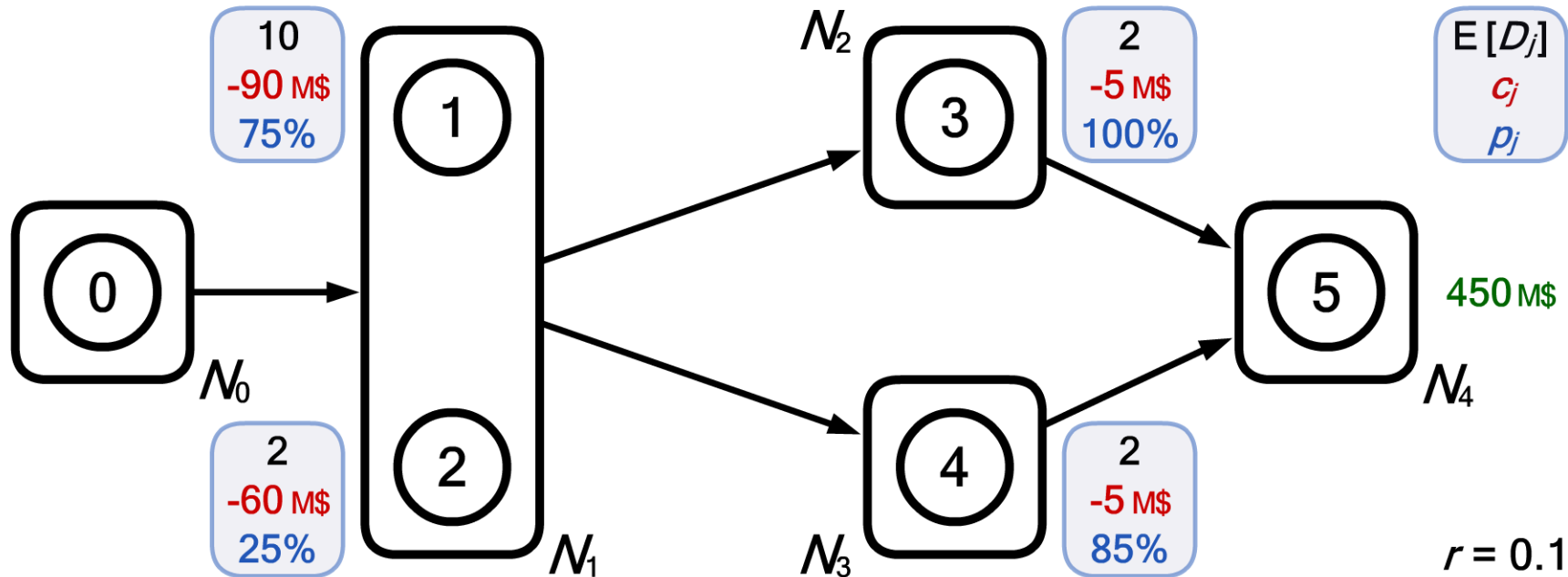
# Deterministic durations



# Deterministic durations



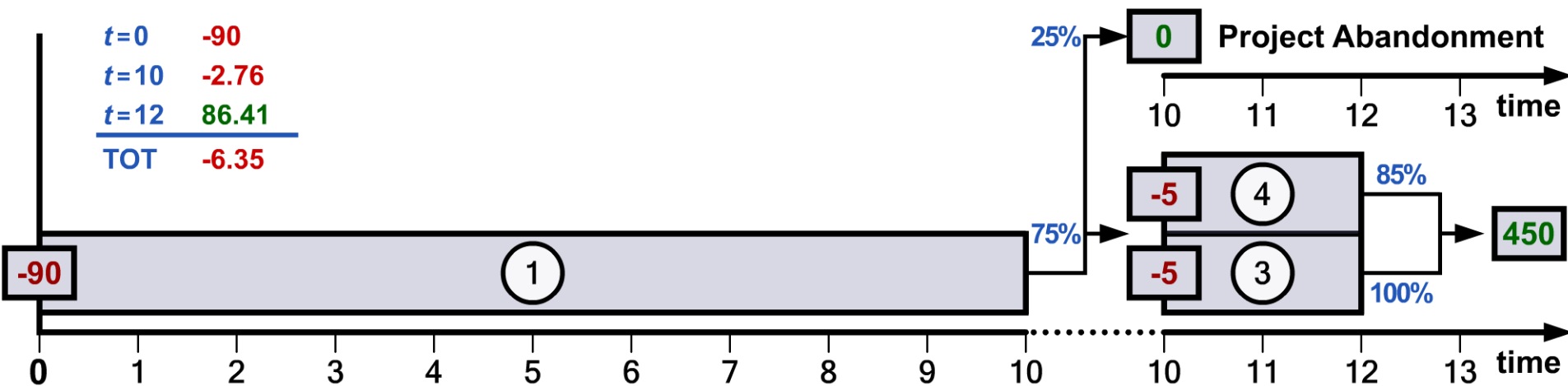
# Deterministic durations



# Deterministic durations

A solution is not a schedule but rather a scheduling policy (even with deterministic durations)

Policy  $\Pi_1$  results in a NPV of **-6.35M\$** if activity durations are deterministic





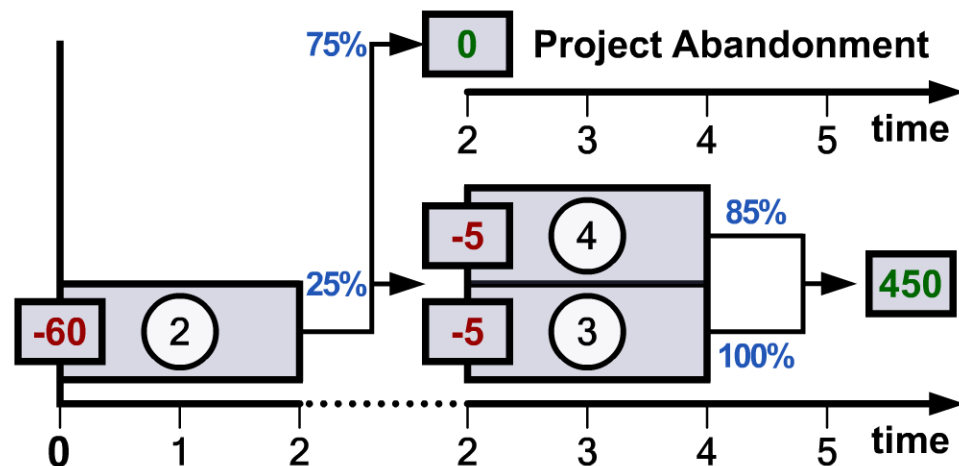
# Deterministic durations

A solution is not a schedule but rather a scheduling policy (even with deterministic durations)

Policy  $\Pi_1$  results in a NPV of **-6.35M\$** if activity durations are deterministic

Policy  $\Pi_2$  is optimal for deterministic durations and yields a NPV of **2.05M\$**

$t=0$	-60
$t=2$	-2.04
$t=4$	64.10
TOT	2.05



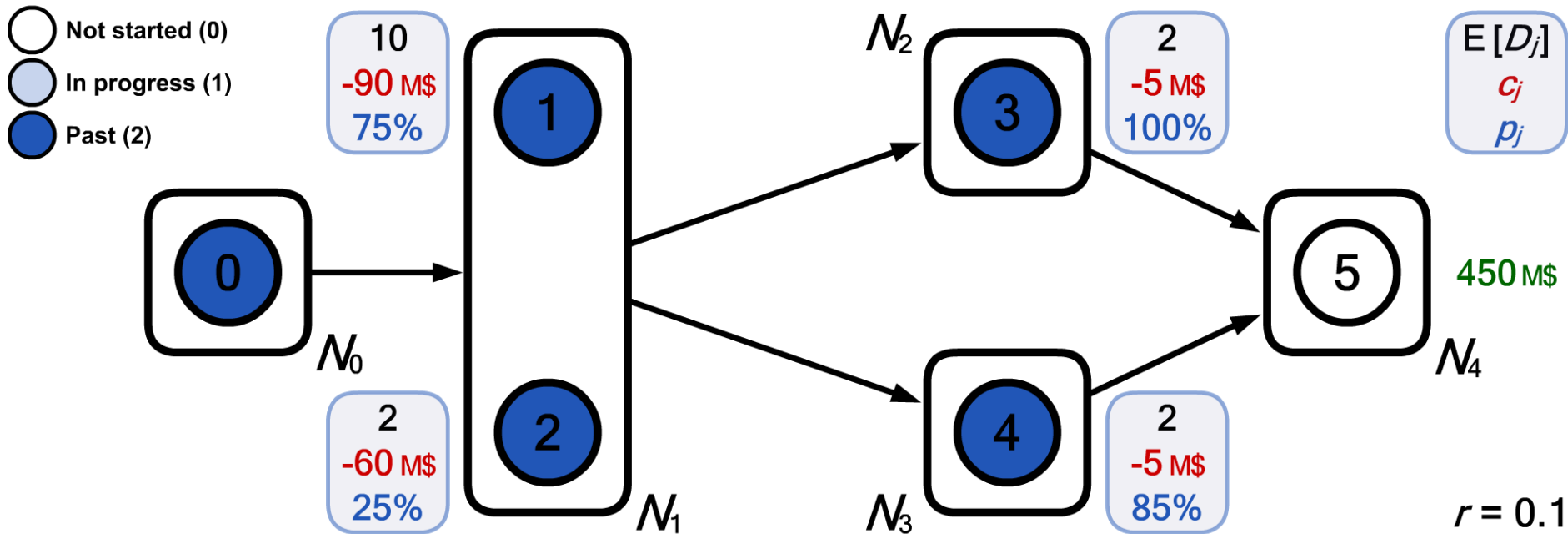
# Stochastic durations: methodology

- Exponentially distributed durations => use of a Continuous-Time Markov Chain (CTMC) to model the statespace
- State of an activity  $j$  at time  $t$  can be:
  - Not started
  - In progress
  - Past (successfully finished, failed or considered redundant because another activity of its module has completed successfully)
- Size of statespace has upper bound  $3^n$ . Most states do not satisfy precedence constraints => a strict definition of the statespace is required. This is studied in Creemers et al. (2010)\*

⇒ Backward SDP-recursion

\*Creemers S, Leus R, Lambrecht M (2010). Scheduling Markovian PERT networks to maximize the net present value. Operations Research Letters, vol. 38, no. 1, pp. 51 - 56.

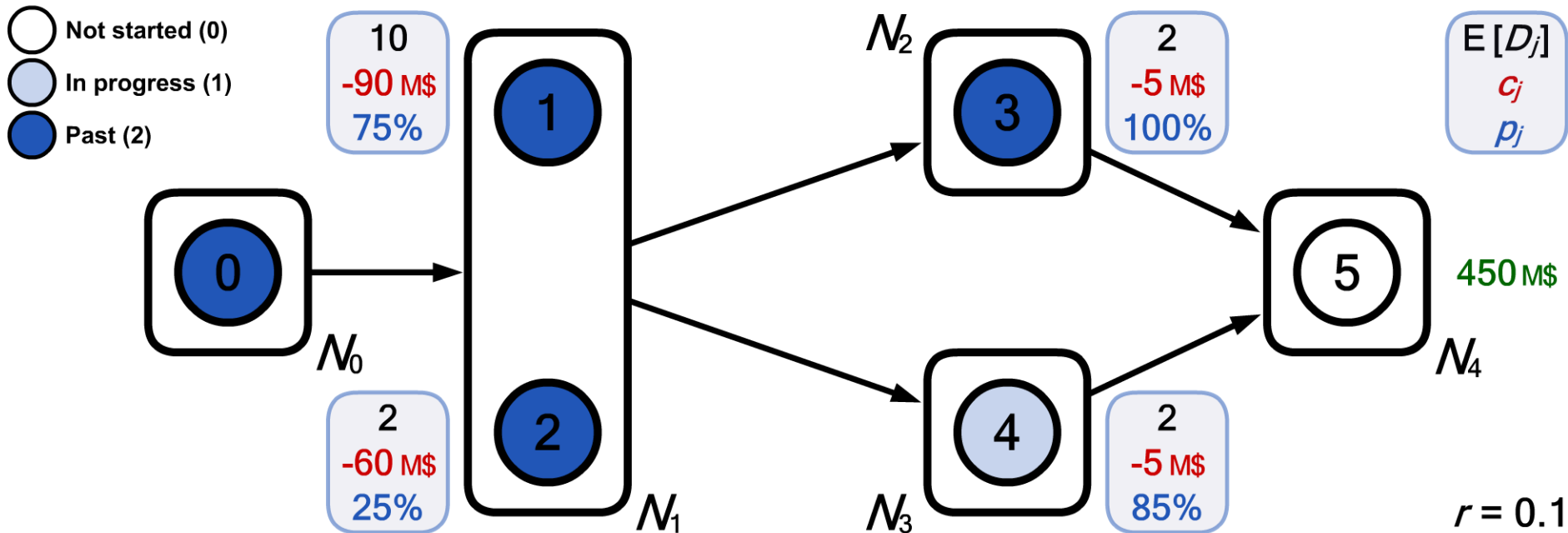
# Stochastic durations



(2,2,2,2,2,0) [450M\$]

Project value upon entry of the final state = project payoff

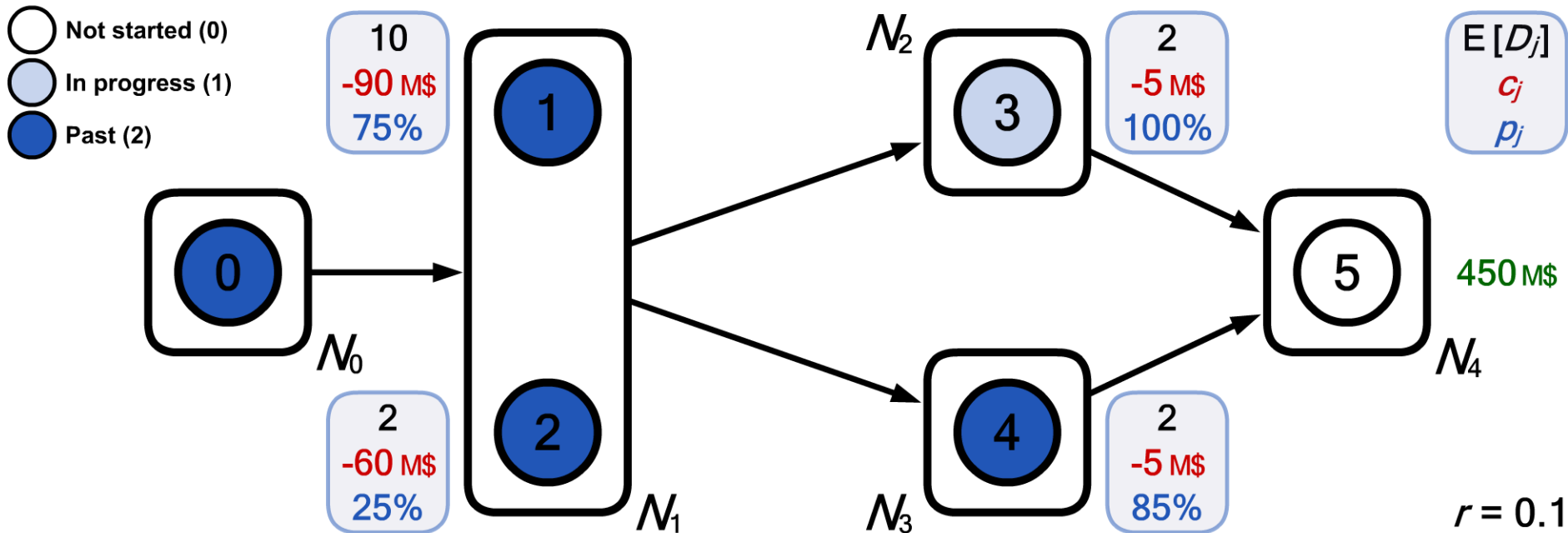
# Stochastic durations



(2,2,2,2,2,0) [450M\$]  
 ↳ (2,2,2,2,1,0) [318.75M\$]

Discount factor:  $(1/E[D_4]) / (r + (1/E[D_4]))$   
 $D_4 = 2 \Rightarrow$  discount factor = 0.83  
 NPV upon state entry if success = 375  
 $p_4 = 0.85 \Rightarrow$  NPV upon state entry = 318.75

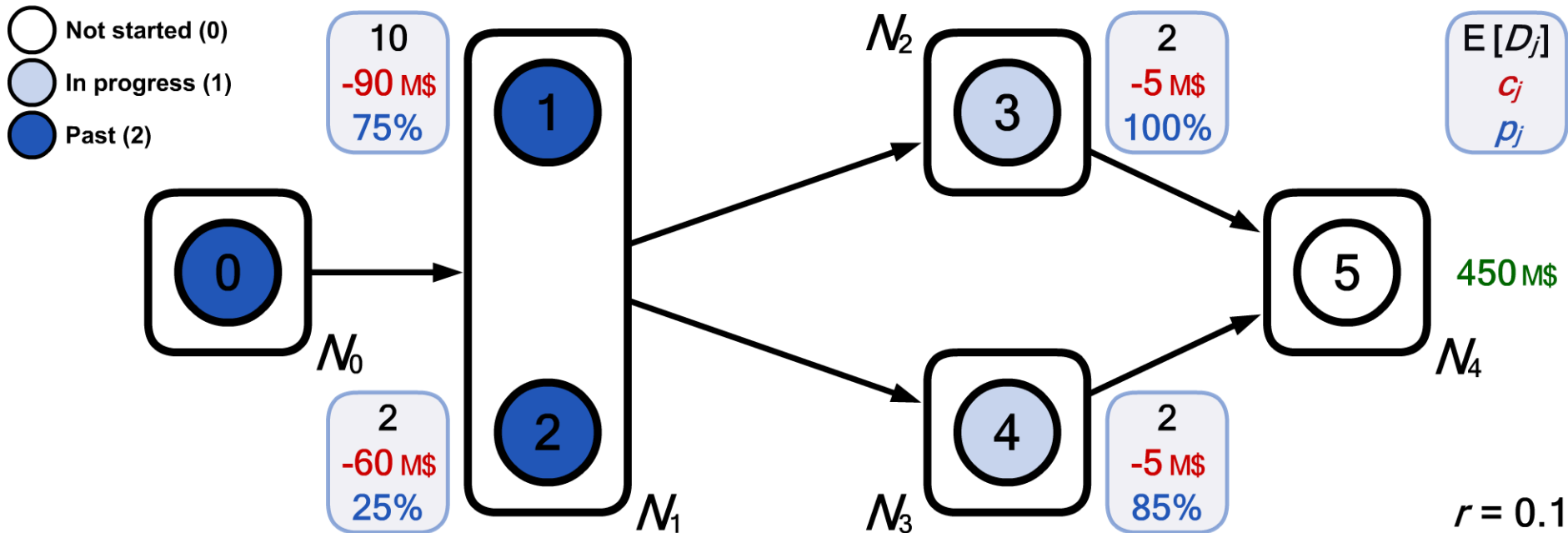
# Stochastic durations



(2,2,2,2,2,0) [450M\$]  
 ↳ (2,2,2,2,1,0) [318.75M\$]  
 ↳ (2,2,2,1,2,0) [375M\$]

Discount factor:  $(1/E[D_3]) / (r + (1/E[D_3]))$   
 $D_3 = 2 \Rightarrow$  discount factor = 0.83  
 NPV upon state entry if success = 375  
 $p_3 = 1.00 \Rightarrow$  NPV upon state entry = 375

# Stochastic durations



(2,2,2,2,2,0) [450M\$]  
 ↳ (2,2,2,2,1,0) [318.75M\$]  
 ↳ (2,2,2,1,2,0) [375M\$]  
 ↳ (2,2,2,1,1,0) [289.77M\$]

Discount factor = 0.91

Probability of finishing activity  $j$  first :  $(1/E[D_j]) / (\sum_k (1/E[D_k]))$   
 => Probability 3 finishes first is 50% and  $p_3 = 100\%$

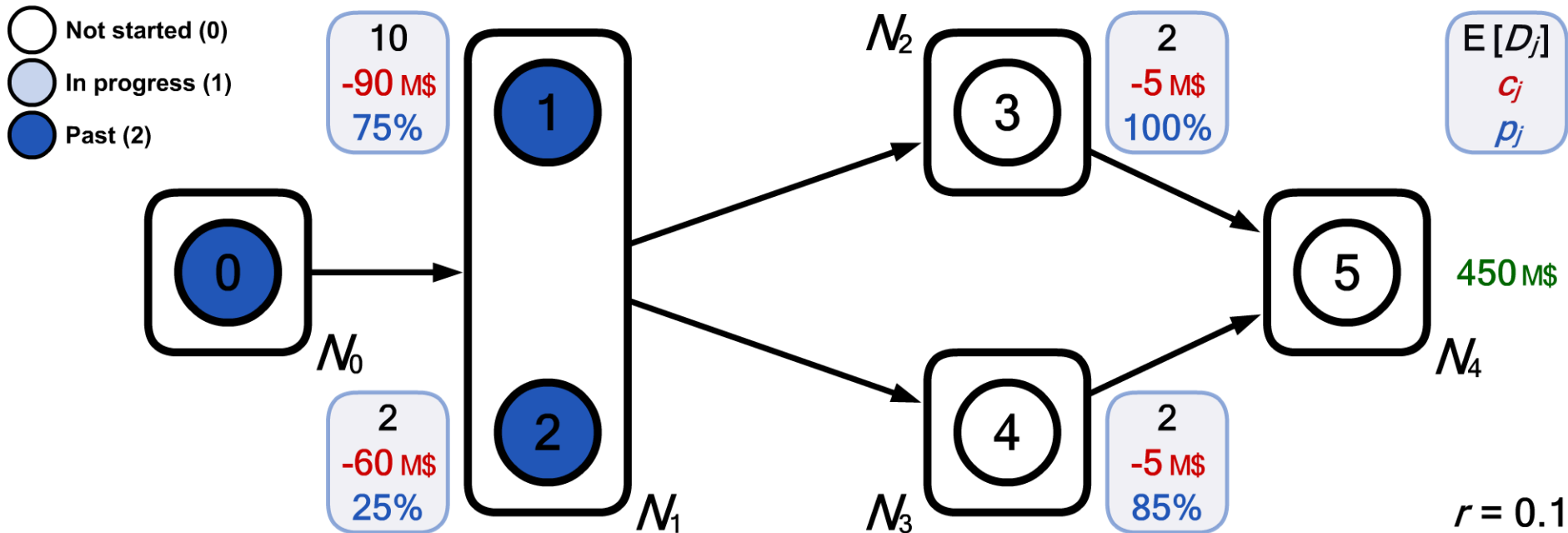
$0.5 \times 0.91 \times 1.00 \times 318.75 = 144.89$

=> Probability 4 finishes first is 50% and  $p_4 = 85\%$

$0.5 \times 0.91 \times 0.85 \times 375 = 144.89$

=> NPV upon state entry = 289.77

# Stochastic durations



(2,2,2,2,2,0) [450M\$]

→ (2,2,2,2,1,0) [318.75M\$]

→ (2,2,2,1,2,0) [375M\$]

→ (2,2,2,1,1,0) [289.77M\$]

→ (2,2,2,0,0,0) [279.77M\$]

3 possible decisions (pick the optimal one):

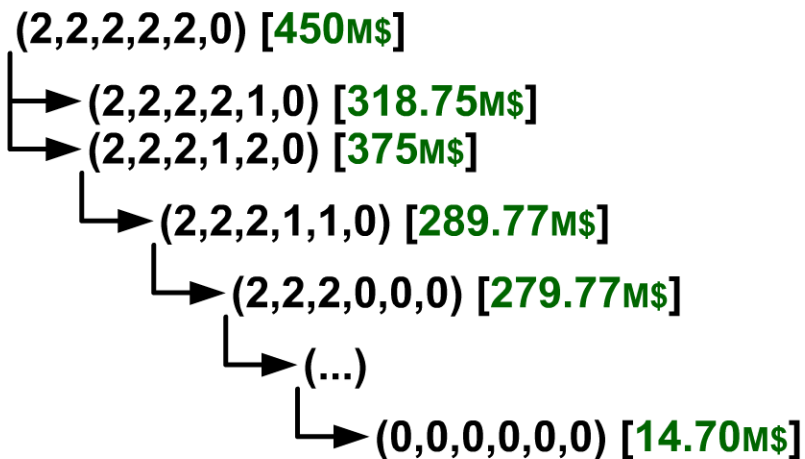
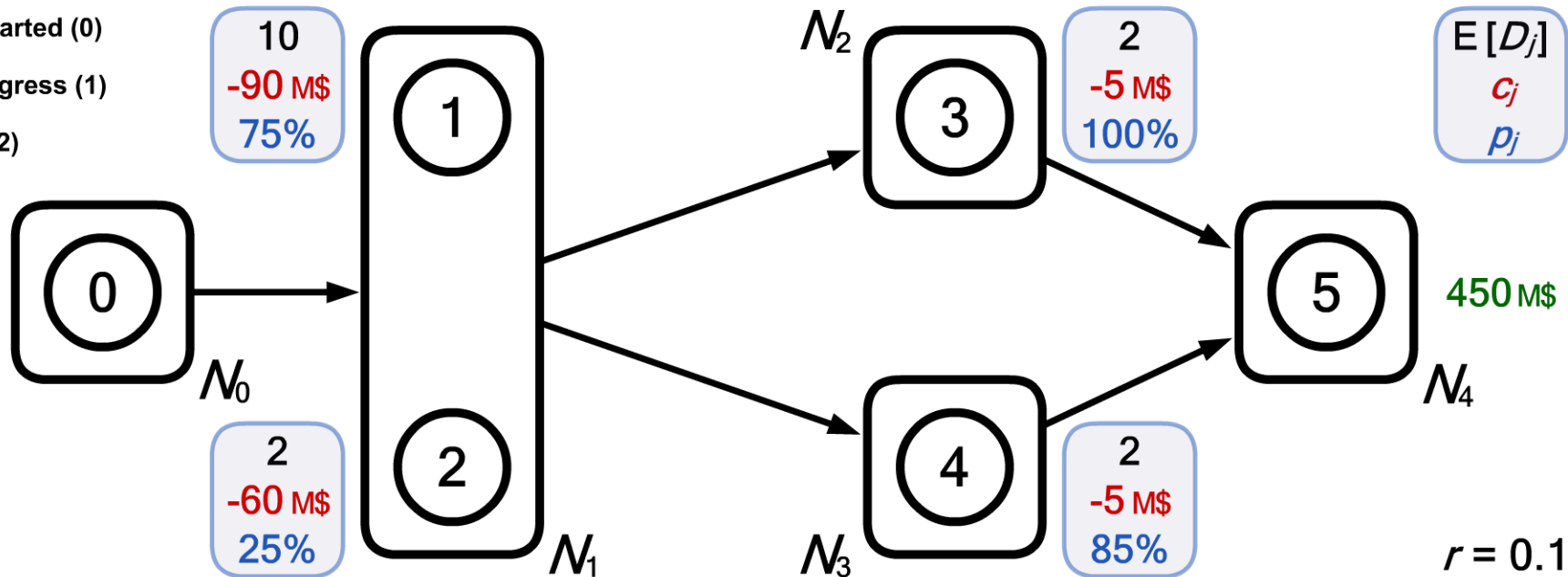
- Start activity 3 => incur cost  $c_3 = -5M\$$   
=> end up in (2,2,2,1,0,0)

- Start activity 4 => incur cost  $c_4 = -5M\$$   
=> end up in (2,2,2,0,1,0)

- Start activity 3 & 4 => incur cost  $c_3 + c_4 = -10M\$$   
=> end up in (2,2,2,1,1,0)[289.77M\$]

# Stochastic durations

- Not started (0)
- In progress (1)
- Past (2)





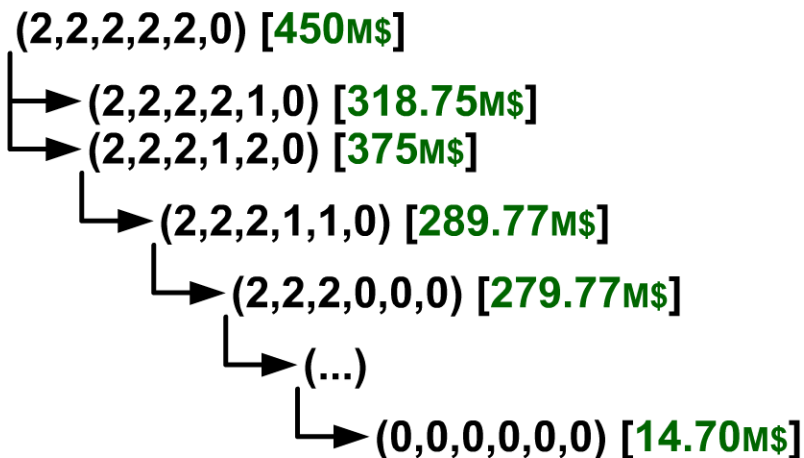
# Stochastic durations

A solution is not a schedule but rather a scheduling policy (even with deterministic durations)

Policy  $\Pi_1$  results in a NPV of **-6.35M\$** if activity durations are deterministic

Policy  $\Pi_2$  is optimal for deterministic durations and yields a NPV of **2.05M\$**

For stochastic durations, policy  $\Pi_1$  is optimal with a NPV of **14.70M\$**



# Results & future work

- Computational results:
  - 1260 randomly generated projects have been solved to optimality

$n$	10	20	30	60	90
<b>CPU (sec)</b>	0.00	0.03	1.95	84.04	4100.52

- Main determinant of computation time = network density (for fixed  $n$ )
- Future work:
  - Using the model to generate insights
  - General activity durations using Phase-type distributions
  - Renewable resources

# Questions?

